

Relative monads and their many guises

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1970 : Devices

In 1970, Walters introduced the notion of **device** to capture classical universal algebra, and which was shown to generalise the notion of monad.

However, the concept was overlooked, and devices were forgotten...

2010: Relative monads

Forty years later, the concept of **relative monad** was introduced by Altenkirch, Chapman, and Uustalu as a generalisation of the concept of monad, from a structured endofunctor to an arbitrary functor with structure.

Crucially, the authors related relative monads to the **relative adjunctions** of Ulmer.

Relative monads and devices

In fact, the definition of relative monad is reminiscent to that of device, and the two turn out to be equivalent.

Relative monads in extension form [ACU10]

Fix a functor $j: A \rightarrow E$.

A j -relative monad comprises

- a functor $t: A \rightarrow E$
- a natural transformation $\eta: j \Rightarrow t$
- a natural transformation

$$\dagger: E(j=, t-) \Rightarrow E(t=, t-)$$

satisfying unit and associativity laws.

Where does the concept of
relative monad
come from, categorically?

Monads as monoids

For any category A , the category $\text{Cat}(A, A)$ of endofunctors on A is equipped with strict monoidal structure given by functor composition.

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Can this be generalised to arbitrary functor categories?

A hint of a solution

ACU show that, assuming the existence of enough (pointwise) extensions along $A \xrightarrow{j} E$, $\text{Cat}(A, E)$ may be equipped with skew-monoidal structure, in which the unit is $A \xrightarrow{j} E$, and the tensor of $A \xrightarrow{f} E$ and $A \xrightarrow{g} E$ is given by left extension:

$$\begin{array}{ccc} A & \xrightarrow{f} & E \\ & \nearrow g & \searrow j \triangleright g \\ A & \xrightarrow{j} & E \end{array}$$

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Furthermore, monoids in this skew-monoidal category of functors are precisely j -relative monads.

However, the assumption of enough left extensions along j is dissatisfying; after all, monads are always monoids in a category of functors.

\mathcal{V} -distributors (AKA \mathcal{V} -profunctors)

Let \mathcal{V} be a monoidal category. A \mathcal{V} -distributor $A \overset{p}{\dashrightarrow} B$ comprises

$$p(b, a) \in \mathcal{V} \quad (b \in B, a \in A)$$

functorial contravariantly in b and covariantly in a .

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\mathcal{V} -distributors $A \overset{p}{\dashrightarrow} B \overset{q}{\dashrightarrow} C$ cannot generally be composed. We will denote by $q \circ p$ the composite when it exists.

\mathcal{V} -forms

Consider \mathcal{V} -distributors

$$A_n \xrightarrow{p_n} A_{n-1} \xrightarrow{p_{n-1}} \dots \xrightarrow{p_2} A_1 \xrightarrow{p_1} A_0$$

A \mathcal{V} -form $p_1, \dots, p_n \Rightarrow q$ comprises a family

$$\left\{ p_1(a_0, a_1) \otimes \dots \otimes p_n(a_{n-1}, a_n) \rightarrow q(a_0, a_n) \right\}_{a_i \in A_i}$$

of morphisms in \mathcal{V} , satisfying naturality laws.

Representable distributors

Let $A \xrightarrow{f} B$ be a \mathcal{V} -functor. There is an induced representable \mathcal{V} -distributor

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A \mathcal{V} -form $B(1, f) \Rightarrow B(1, g)$ is just a \mathcal{V} -natural transformation $f \Rightarrow g$.

Multicategories of endofunctors

There is a submulticategory $\mathcal{V}\text{-Cat}[A, A]$ of $\mathcal{V}\text{-Dist}[A, A]$ formed by the representable \mathcal{V} -distributors.

We may always compose representable \mathcal{V} -distributors, and so $\mathcal{V}\text{-Cat}[A, A]$ is representable: i.e. a monoidal category. It is equivalent to the usual strict monoidal category $\mathcal{V}\text{-Cat}(A, A)$ of \mathcal{V} -endofunctors on A .

Skew composition I

Let A and E be categories. Given functors

$$A \xrightarrow{f} E \quad A \xrightarrow{g} E$$

We cannot in general form any sort of composite $A \xrightarrow{g \cdot f} E$, since the codomain of f does not match the domain of g .

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However, suppose we had a fixed functor $E \xrightarrow{j^*} A$. Then we could form a composite

$$A \xrightarrow{f} E \xrightarrow{j^*} A \xrightarrow{g} E$$

Skew composition II

Fix a \mathcal{V} -functor $A \xrightarrow{j} E$. Given \mathcal{V} -functors

$$A \xrightarrow{f} E \quad A \xrightarrow{g} E$$

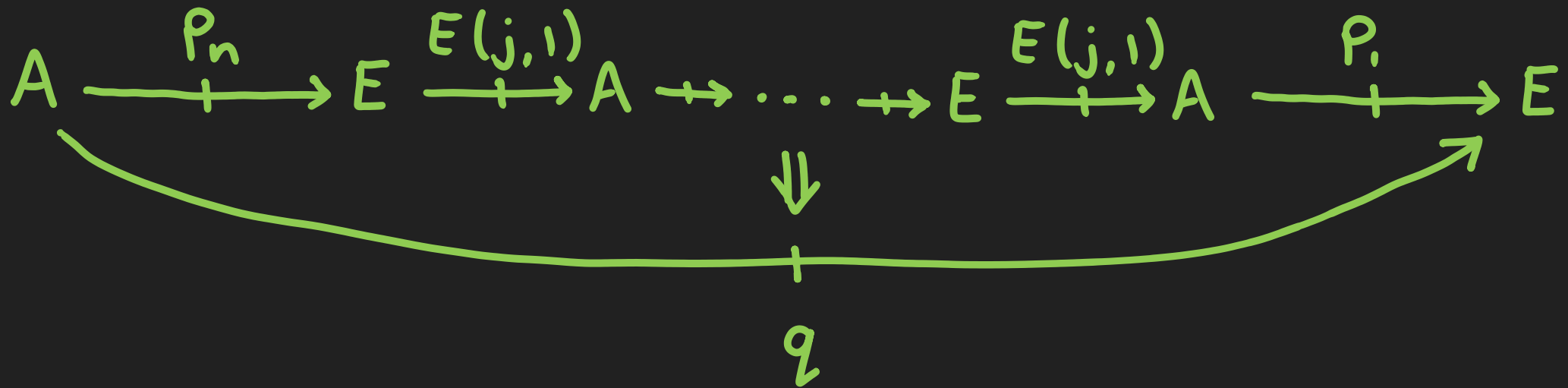
we may form a chain

$$A \xrightarrow{E(1,f)} E \xrightarrow{E(j,1)} A \xrightarrow{E(1,g)} E$$

by considering f, g as representable \mathcal{V} -distributors,
and using the corepresentable $E(j,1)$ to facilitate
composition.

Skew-multicategories of distributors

Theorem. Let $j: A \rightarrow E$ be a \mathcal{V} -functor.
There is a skew-multicategory $\mathcal{V}\text{-Dist}[j]$
whose objects are \mathcal{V} -distributors $A \leftrightarrow E$
and whose multimorphisms are \mathcal{V} -forms



Skew-multicategories of functors

There is a sub skew-multicategory $\mathcal{V}\text{-Cat}[j]$ of $\mathcal{V}\text{-Dist}[j]$ formed by the representables.

$$\begin{array}{ccccccc} A & \xrightarrow{E(1, f_n)} & E & \xrightarrow{E(j, 1)} & A & \rightsquigarrow \dots \rightsquigarrow & E & \xrightarrow{E(j, 1)} & A & \xrightarrow{E(1, f_1)} & E \\ & & & & & \downarrow & & & & & \nearrow \\ & & & & & E(1, g) & & & & & \end{array}$$

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We cannot generally compose representable \mathcal{V} -distributors with corepresentable ones, so $\mathcal{V}\text{-Cat}[j]$ is usually not representable.

Monoids in $\mathcal{V}\text{-Cat}[j]$

A monoid in $\mathcal{V}\text{-Cat}[j]$ comprises

- a \mathcal{V} -functor $t: A \rightarrow E$
- a \mathcal{V} -natural transformation $\eta: j \Rightarrow t$
- a \mathcal{V} -form $\mu: E(1, t), E(j, t) \Rightarrow E(1, t)$

satisfying unit and associativity laws.

This looks quite suggestive.

Transposition

Suppose that \mathcal{V} -distributors may be composed. Then there is an adjunction in $\mathcal{V}\text{-Dist}$

$$E(1, t) \dashv E(t, 1)$$

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$$\underline{E(1, t)} \odot E(j, t) \Rightarrow E(1, t)$$

are, by transposition, in bijection with \mathcal{V} -forms

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A similar principle applies in the absence of composites.

Relative monads as monoids

Theorem. Let $A \xrightarrow{j} E$ be a \mathcal{V} -functor.

There is an isomorphism of categories

$$\text{Mon}(\mathcal{V}\text{-cat}[j]) \cong \text{RMnd}(j)$$

Hence, just as every monad is a monoid in a monoidal category, every relative monad is a monoid in a skew-multicategory.

Note that we do not need to impose any assumptions on j .

Equivalent concepts

Relative monads

Monoids in $\mathcal{V}\text{-Cat}[j]$

Representability of $\mathcal{V}\text{-Cat}[j]$

It is natural to ask when the skew-multicategory $\mathcal{V}\text{-Cat}[j]$ is representable by a skew-monoidal category.

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Theorem. Suppose that left extensions of \mathcal{V} -functors $A \rightarrow E$ along $A \xrightarrow{j} E$ are admitted. Then $\mathcal{V}\text{-Cat}[j]$ is representable.

We thereby recover ACU's characterisation of relative monads as monoids in a SkMC.

Equivalent concepts

Relative monads

Monoids in $\mathcal{V}\text{-Cat}[j]$

SkMulticat

SkMonCat

Formal mw-monads

In 2014 Lack-Street carried out a similar study of skew-monoidal hom-categories to study the extension form of monads. They defined a notion of **formal mw-monad**, noting an apparent similarity to the relative monads of ACU.

Monoids in \mathcal{V} -Dist[j]

We can recover j -relative monads from monoids in \mathcal{V} -Dist.

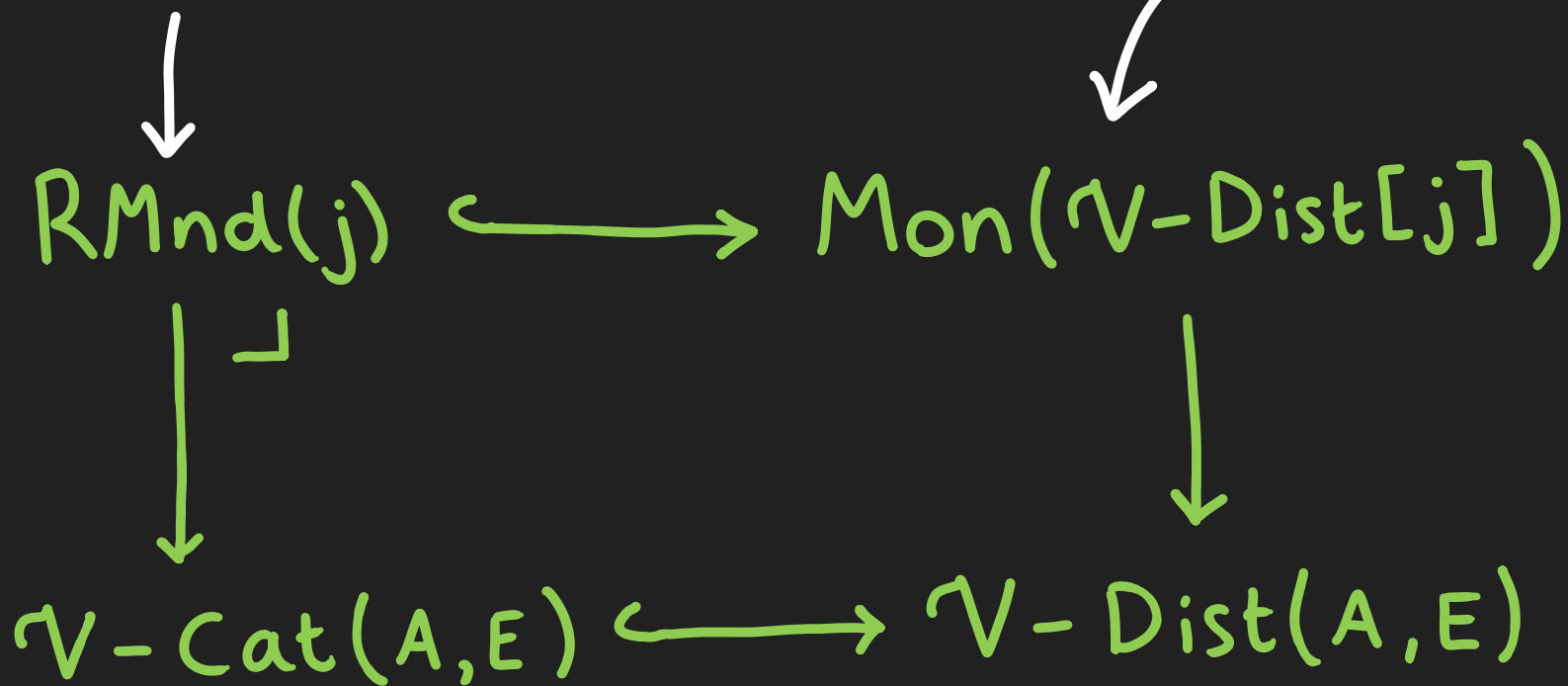
Theorem. There is a pullback

$$\begin{array}{ccc} \text{Mon}(\mathcal{V}\text{-Cat}[j]) & \longleftrightarrow & \text{Mon}(\mathcal{V}\text{-Dist}[j]) \\ \downarrow \lrcorner & & \downarrow \\ \mathcal{V}\text{-Cat}(A, E) & \longleftrightarrow & \mathcal{V}\text{-Dist}(A, E) \end{array}$$

Representable formal mw-monads

representable
formal mw-monads
[ACU10]

formal mw-monads
in \mathcal{V} -Dist [LS14]



Equivalent concepts

Relative monads

Monoids in $\mathcal{V}\text{-Cat}[j]$

SkMulticat

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Representable monoids
in $\mathcal{V}\text{-Dist}[j]$

Relative monads as monads in Dist

In 1975, Diers introduced the notion of j -monad, for $A \xrightarrow{j} E$ a dense, fully faithful functor.

A j -monad is a monoid in $\text{Dist}(A, A)$ whose underlying endodistributor is of the form

$$E(j, t) : A \rightrightarrows A$$

for some functor $A \xrightarrow{t} E$.

In 2016, Lucyshyn-Wright also studied such monoids, calling such distributors **copresheaf-representable**.

Distributors to endodistributors

Theorem. There is a skew-multifunctor

$$\mathcal{V}\text{-Cat}[j] \rightarrow \mathcal{V}\text{-Dist}[A, A]$$

$$(A \xrightarrow{t} E) \mapsto E(j, t) : A \dashrightarrow A$$

which is fully faithful when j is dense.

Hence, when j is dense, we have a pullback

$$\begin{array}{ccc} \text{Mon}(\mathcal{V}\text{-Cat}[j]) & \hookrightarrow & \text{Mon}(\mathcal{V}\text{-Dist}[A, A]) \\ \downarrow \lrcorner & & \downarrow \\ \mathcal{V}\text{-Cat}(A, E) & \xrightarrow{E(j, -)} & \mathcal{V}\text{-Dist}(A, A) \end{array}$$

Equivalent concepts

Relative monads

Monoids in $\mathcal{V}\text{-Cat}[j]$

SkMulticat

SkMonCat

Representable monoids
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j -representable monoids
in $\mathcal{V}\text{-Dist}[A, A]$ *

Monoidality

Theorem. Suppose that left extensions of \mathcal{V} -functors $A \rightarrow E$ along $A \overset{j}{\rightarrow} E$ are admitted. Then $\mathcal{V}\text{-Cat}[j]$ is representable.

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Theorem. Suppose that left extensions of \mathcal{V} -functors $A \rightarrow E$ along $A \xrightarrow{j} E$ are admitted. Then $\mathcal{V}\text{-Cat}[j]$ is representable.

Furthermore if ... then $\mathcal{V}\text{-Cat}[j]$ is

j is dense		left-normal
j is f.f.		right-normal
left extensions along j are j -absolute		associative-normal

A characterisation of free cocompletions

Theorem. A \mathcal{V} -functor $A \xrightarrow{j} E$ exhibits the free cocompletion of A under a class Φ of weights if and only if

- j is dense and fully faithful
- left extensions along j exist and are j -absolute

Cocontinuous monads

Corollary. Let Φ be a class of weights and let \mathcal{A} be a \mathcal{V} -category admitting a free Φ -cocompletion $\varphi_{\mathcal{A}}: \mathcal{A} \rightarrow \Phi\mathcal{A}$. Then

$$\mathbf{RMnd}(\varphi_{\mathcal{A}}) \simeq \mathbf{Mnd}_{\Phi}(\Phi\mathcal{A})$$

$\varphi_{\mathcal{A}}$ -relative monads are equivalent to Φ -cocontinuous \mathcal{V} -monads on $\Phi\mathcal{A}$.

Cocontinuous monads

Corollary. Let Φ be a class of weights and let A be a \mathcal{V} -category admitting a free Φ -cocompletion $\varphi_A: A \rightarrow \Phi A$. Then

$$\mathbf{RMnd}(\varphi_A) \simeq \mathbf{Mnd}_{\Phi}(\Phi A)$$

φ_A -relative monads are equivalent to Φ -cocontinuous \mathcal{V} -monads on ΦA .

Example. Finitary monads on \mathbf{Set} are $(\mathbf{Set}_f \hookrightarrow \mathbf{Set})$ -relative monads.

Equivalent concepts

Relative monads

Monoids in $\mathcal{V}\text{-Cat}[j]$

SkMulticat

SkMonCat

Representable monoids
in $\mathcal{V}\text{-Dist}[j]$

j -cocontinuous
 \mathcal{V} -monads *

j -representable monoids
in $\mathcal{V}\text{-Dist}[A, A]$ *

A formal theory of relative monads

For the purpose of this talk, I have worked in the setting of categories enriched in a monoidal category.

However, the results hold much more generally: in any virtual equipment.

In particular, this includes $\mathcal{V}\text{-Cat}$ for any virtual double category \mathcal{V} with restrictions.

Summary

- Relative monads have been rediscovered many times in the last 50 years, albeit in different guises.
- We may understand these fruitfully through distributors and 'skew composition'.
- These viewpoints are useful for proving general theorems about relative monads.

Equivalent concepts

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j -cocontinuous
 \mathcal{V} -monads *

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And others...

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